## 6 <br> COMPUTATIONAL RESULTS

Computational tests were perfomed to analyze the performance of the proposed reformulation schemes and solution algorithm. All testes were conducted on a computer with processor Pentium $4,3.00 \mathrm{GHz}$ and 2 GB of RAM. The models and algorithms were implemented using the modeling language MOSEL and solved by XPRESS 19.00.04.

The first results are those obtained for the set of instances described in Viswanath et al. [66]. These are all small-size problems which served as a "proof of correctness" for the proposed methodology. Since no other work in the literature deals with the problem in its original form (remember that [43] dismisses the probabilistic nature of the problem by assuming that investment on an edge completely eliminates the probability of that edge failing afterwards), several other instances were created in order to assess the performance of the methodology for medium and large-size instances of the problem.

The remainder of this Chapter is organized as follows: Section (6.1) presents the results for the instances provided in [66], Section (6.2) describes how the medium and large-size instances were generated and presents results for the former while Section (6.3) discusses the results for the latter.

## 6.1

Instances from the literature

All the instances solved in [66] refer to a graph which contains 4 vertices and 5 edges, as depicted in Figure 6-1. There is a total of 28 instances which are detailed in Table 6.1: they differ from each other in the investment and transportation costs associated with each edge (columns InvCost and TranspCost, respectively), maximum budget (column Budget), penalty for not fulfilling the demand associated to a vertex (column Penalty) and initial and final survival probabilities (initial survival probability is equal to $70 \%$ for all edges in instances 1 through 14 and equal to $60 \%$ in instances 15 through 28 and column

SurvProbInv). Vertices $O$ and $D$ are the origin and destination for a unit commodity that must flow through the network.


Figure 6-1 - Graph corresponding to the instances solved in Viswanath et al. [66]

For all these instances, there are $32\left(2^{5}\right)$ scenarios of network configuration - given by all the possible combinations of the availability of the edges - and the first step of the proposed methodology determines that the minimum cost network flow problem corresponding to each one of these configurations must be solved independently. For this set of instances, total solution time of the network flow problems for all scenarios is minuscule.

Once these optimal values are known, they are used as coefficients in the objective function of the main problem, which is then solved by the algorithm outlined in Chapter 4. All instances were solved to optimality in less than 1.0 second and average solution time was 0.313 second. Details are provided in Table 6.2 where the column Id indicates the instance identification, column OptVal presents the value of the optimal solution, column \# Iter indicates the number of iterations of the algorithm until convergence was achieved and column TotalTime the time it took for the algorithm to complete.

Table 6-1 - Description of the instances provided in Viswanath et al. [66]

| Id | SurvProbInv | InvCost | Budget | TranspCost |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $1 / 15$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,1\}$ | 2 | $\{10,10,10,10,10\}$ | Penalty |
| $2 / 16$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,1\}$ | 3 | $\{10,10,10,10,10\}$ | 31 |
| $3 / 17$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,1\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $4 / 18$ | $\{80 \%, 80 \%, 90 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,1\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $5 / 19$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{2,1,1,1,1\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $6 / 20$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,2,1,1,1\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $7 / 21$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,2,1,1\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $8 / 22$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,2,1\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $9 / 23$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,2\}$ | 3 | $\{10,10,15,30,10\}$ | 41 |
| $10 / 24$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{2,1,1,1,1\}$ | 3 | $\{10,20,10,15,10\}$ | 31 |
| $11 / 25$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{2,1,1,1,1\}$ | 3 | $\{10,20,10,15,10\}$ | 43.9 |
| $12 / 26$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{2,1,1,1,1\}$ | 3 | $\{10,20,10,15,10\}$ | 57.3 |
| $13 / 27$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,1\}$ | 3 | $\{10,15,5,15,10\}$ | 26 |
| $14 / 28$ | $\{80 \%, 80 \%, 80 \%, 80 \%, 80 \%\}$ | $\{1,1,1,1,1\}$ | 3 | $\{10,15,1,15,10\}$ | 2 |

Table 6-2 - Results of the instances provided in Viswanath et al. [66]

| Id | OptVal | \# Iter | TotalTime |
| :---: | :---: | :---: | :---: |
| 1 | 21.9961 | 3 | 0.137 |
| 2 | 21.7155 | 2 | 0.082 |
| 3 | 26.8835 | 3 | 0.123 |
| 4 | 26.8494 | 4 | 0.225 |
| 5 | 26.9087 | 3 | 0.120 |
| 6 | 26.9681 | 3 | 0.106 |
| 7 | 26.8835 | 3 | 0.136 |
| 8 | 26.8835 | 3 | 0.122 |
| 9 | 26.9681 | 3 | 0.129 |
| 10 | 26.9601 | 4 | 0.257 |
| 11 | 29.0251 | 3 | 0.144 |
| 12 | 31.0963 | 3 | 0.297 |
| 13 | 25.1315 | 5 | 1.000 |
| 14 | 23.0995 | 3 | 0.359 |
| 15 | 22.5114 | 3 | 0.302 |
| 16 | 22.0285 | 2 | 0.187 |
| 17 | 26.9725 | 3 | 0.359 |
| 18 | 26.9638 | 3 | 0.531 |
| 19 | 27.0157 | 3 | 0.359 |
| 20 | 27.1194 | 3 | 0.375 |
| 21 | 26.9725 | 3 | 0.360 |
| 22 | 26.9725 | 3 | 0.421 |
| 23 | 27.1194 | 3 | 0.359 |
| 24 | 27.0074 | 3 | 0.422 |
| 25 | 28.8943 | 3 | 0.421 |
| 26 | 32.0447 | 3 | 0.375 |
| 27 | 25.1565 | 3 | 0.625 |
| 28 | 23.1405 | 2 | 0.235 |

## 6.2

## Medium-size instances

Given the lack of additional instances of the problem available in the literature, we developed an instance generator which was then used to test the proposed methodology.

The instances were created by randomly selecting the location of a given number of vertices within a region defined by minimum and maximum values for the $x$ and $y$ coordinates. Next, a predefined number of edges connecting the vertices was created (the resulting graph was checked for connectedness in order to avoid trivial and meaningless solutions) and the Euclidean distance between the corresponding vertices was assigned as the transportation cost of each edge. Preand post-investment survival probabilities were assigned to each edge and, for the large instances presented in Section 6.3, scenarios of network configuration were generated based on the initial survival probability of each edge.

Next, in Tables 6.3 and 6.4, we present the results for a total of 30 instances which were all solved by the algorithm designed in Chapter 4 (ALG2) with full scenario enumeration and tolerance level set to no more than $1 \%$. The table provides the following information: column Id identifies the instance, columns \# Vertices and \# Edges indicates, respectively, the number of vertices and edges of the graph, column \# Scen provides the number of scenarios of network configuration used in each problem; column $U B$ reports the value of the best solution found while column $L B$ indicates the value of the solution to the last approximated problem (i.e., the one which is solved by considering the set of cuts that approximate the exponential function), column \% Gap presents the percentage gap between the upper and lower bounds and column ErrTol contains the maximum acceptable error, which is the stopping criterium for the algorithm; column \# Iter indicates the number of iterations of the algorithm needed to reach the final solution, and column MainTime report the total time for the convergence of the algorithm (the time needed for the solution of the independent scenariospecific network flow problems is not reported but they are usually orders of magnitude smaller than the time it takes for the algorithm to converge which thus represents the bottleneck of the methodology).

Table 6-3-Results for the medium-size instances

| Id | \# Vertices | \# Edges | \# Scen | UB | LB | \% Gap | ErrTol | \# Iter | MainTime |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| v5e6A | 5 | 6 | 64 | 219.423 | 219.423 | $0.00 \%$ | $0.1 \%$ | 5 | 0.437 |
| v5e6B | 5 | 6 | 64 | 50.6467 | 50.6467 | $0.00 \%$ | $0.1 \%$ | 3 | 0.140 |
| v5e6C | 5 | 6 | 64 | 143.923 | 143.923 | $0.00 \%$ | $0.1 \%$ | 4 | 0.281 |
| v5e6D | 5 | 6 | 64 | 107.918 | 107.918 | $0.00 \%$ | $0.1 \%$ | 4 | 0.250 |
| v5e6E | 5 | 6 | 64 | 277.904 | 277.904 | $0.00 \%$ | $0.1 \%$ | 4 | 0.281 |
| v6e8A | 6 | 8 | 256 | 361.265 | 361.265 | $0.00 \%$ | $0.1 \%$ | 6 | 1.796 |
| v6e8B | 6 | 8 | 256 | 45.9159 | 45.9159 | $0.00 \%$ | $0.1 \%$ | 5 | 1.328 |
| v6e8C | 6 | 8 | 256 | 350.268 | 350.268 | $0.00 \%$ | $0.1 \%$ | 5 | 1.594 |
| v6e8D | 6 | 8 | 256 | 110.915 | 110.915 | $0.00 \%$ | $0.1 \%$ | 5 | 1.062 |
| v6e8E | 6 | 8 | 256 | 65.4201 | 65.4201 | $0.00 \%$ | $0.1 \%$ | 3 | 0.328 |
| v7e10A | 7 | 10 | 1024 | 122.857 | 122.851 | $0.0049 \%$ | $0.1 \%$ | 4 | 3.735 |
| v7e10B | 7 | 10 | 1024 | 201.934 | 201.926 | $0.0040 \%$ | $0.1 \%$ | 5 | 4.468 |
| v7e10C | 7 | 10 | 1024 | 104.863 | 104.857 | $0.0057 \%$ | $0.1 \%$ | 4 | 3.063 |
| v7e10D | 7 | 10 | 1024 | 158.868 | 158.861 | $0.0044 \%$ | $0.1 \%$ | 5 | 4.063 |
| v7e10E | 7 | 10 | 1024 | 75.5659 | 75.5619 | $0.0053 \%$ | $0.1 \%$ | 5 | 5.704 |

Table 6-4 - Results for the medium-size instances

| Id | \# Vertices | \# Edges | \# Scen | $\boldsymbol{U B}$ | $\boldsymbol{L B}$ | \% Gap | ErrTol | \# Iter | MainTime |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| v7e11A | 7 | 11 | 2048 | 306.692 | 306.672 | $0.0065 \%$ | $0.1 \%$ | 8 | 79.143 |
| v7e11B | 7 | 11 | 2048 | 252.026 | 251.985 | $0.0163 \%$ | $0.1 \%$ | 6 | 39.344 |
| v7e11C | 7 | 11 | 2048 | 66.7778 | 66.7561 | $0.0325 \%$ | $0.1 \%$ | 5 | 15.593 |
| v7e11D | 7 | 11 | 2048 | 312.985 | 312.959 | $0.0083 \%$ | $0.1 \%$ | 6 | 38.875 |
| v7e11E | 7 | 11 | 2048 | 58.0173 | 57.9977 | $0.0338 \%$ | $0.1 \%$ | 5 | 27.859 |
| v8e12A | 8 | 12 | 4096 | 31.9086 | 31.5927 | $0.99 \%$ | $1 \%$ | 11 | 405.395 |
| v8e12B | 8 | 12 | 4096 | 141.750 | 141.515 | $0.1658 \%$ | $1 \%$ | 6 | 242.878 |
| v8e12C | 8 | 12 | 4096 | 97.1507 | 97.0042 | $0.1508 \%$ | $1 \%$ | 4 | 40.219 |
| v8e12D | 8 | 12 | 4096 | 49.6668 | 49.4192 | $0.4985 \%$ | $1 \%$ | 7 | 286.253 |
| v8e12E | 8 | 12 | 4096 | 155.492 | 155.381 | $0.0714 \%$ | $1 \%$ | 4 | 71.297 |
| v8e12b6A | 8 | 12 | 4096 | 40.8165 | 40.6271 | $0.4640 \%$ | $1 \%$ | 6 | 247.645 |
| v8e12b6B | 8 | 12 | 4096 | 28.5302 | 28.2736 | $0.8994 \%$ | $1 \%$ | 8 | 831.795 |
| v8e12b6C | 8 | 12 | 4096 | 22.3931 | 22.1859 | $0.9253 \%$ | $1 \%$ | 7 | 449.365 |
| v8e12b6D | 8 | 12 | 4096 | 66.6392 | 66.3808 | $0.3878 \%$ | $1 \%$ | 8 | 992.796 |
| v8e12b6E | 8 | 12 | 4096 | 105.355 | 104.637 | $0.6815 \%$ | $1 \%$ | 5 | 113.814 |

## 6.3

## Large-size instances

The set of instances in Section 6.2 involved graphs with a maximum of twelve edges and 4096 possible scenarios of network configuration. The total time required to solve these problems clearly shows how the computational effort increased very rapidly with respect to the number of edges - just as an illustration of this fact, the average time needed to solve the instances with 11 edges was 40.2 seconds, while the average time consumed by the algorithm in solving the instances with 12 edges was 368.1 seconds.

A critical example is provided by an instance of the problem with 10 vertices and 15 edges (and, consequently, 32768 possible scenarios of network configuration) which was solved by full scenario enumeration. Figure 6-2 below presents the performance of the algorithm - data points represent the upper and lower bounds obtained at each iteration:


Figure 6-2 - Algorithm perfomance on an 15-edge instance with full scenario enumeration

While the previous instances converged to solutions with gaps not larger than $1 \%$ after no more than 17 minutes, in the case of the 15 -edge instance it took a total of 25 hours for the algorithm to narrow the gap down to $2.57 \%$. This clearly leads to the conclusion that full scenario enumeration is currently not a
viable option when one tries to solve large scale problems and a sample-based version of the problem - such as that suggested in Chapter 5 - becomes a necessity.

In Table 6.5, we present results for 18 instances with the number of edges ranging from 15 to 40 , also constructed according to the description given in the previous Section. All instances were solved to a maximum gap of $0.87 \%$. Compared to Table 6.4, there is an additional column \# TotScen where the number of possible scenarios of network configuration is reported - column \# Scen indicates the number of scenarios actually used when solving the problem.

The instances with 15 edges (v10e15_1, v10e15_2 and v10e15_3) all refer to the same graph of the example for which the convergence of the algorithm was shown in Figure 6-2. Each one of them was solved using a different set of 500 scenarios (out of the 32768 possible network configurations), sampled according to the initial probability distribution of the edges' availabilities. It is interesting to observe that even though the number of scenarios used in these instances is significantly smaller than the total number of possible scenarios, the solutions found for these problems in under 60 seconds have an objective function value which is close to that found after 25 hours in the case of full scenario enumeration.

A significant increase in computational times was observed when solving the instances with 20 edges (v10e20_1 through v10e20_5) and 500 scenarios. All instances refer to the same graph and were solved using different sets of scenarios to a maximum gap of $0.868 \%$, including one instance which was solved to optimality. The difference between the minimum and maximum optimal values obtained for all five problems was $5.49 \%$, which seems like a reasonable compromise considering that the number of scenarios actually used to solve the problems represents a very small fraction (namely, $0.048 \%$ ) of all possible network configurations.

As the number of binary variables increase so does the computational effort required to solve the problems at each iteration, leading to a compromise between network size (which dictates the number of binary variables) and number of sampled scenarios. This was done when solving the instances with 25,30 and 40 edges: the number of scenarios utilized was reduced to 300 in the former and 200 in the last two cases.

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Table 6-5 - Results for the large-size instances

| Id | \# Vertices | \# Edges | \# Scen | \# TotScen | UB | LB | \% Gap | ErrTol | \# Iter | MainTime |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| v10e15_1 | 10 | 15 | 500 | $3.28 \mathrm{E}+4$ | 44.5735 | 44.4068 | $0.374 \%$ | $1 \%$ | 6 | 41.032 |
| v10e15_2 | 10 | 15 | 500 | $3.28 \mathrm{E}+4$ | 44.8405 | 44.8405 | $0 \%$ | $1 \%$ | 6 | 27.562 |
| v10e15_3 | 10 | 15 | 500 | $3.28 \mathrm{E}+4$ | 47.0452 | 46.9256 | $0.254 \%$ | $1 \%$ | 6 | 27.578 |
| v10e20_1 | 10 | 20 | 500 | $1.05 \mathrm{E}+6$ | 81.8326 | 81.7279 | $0.128 \%$ | $1 \%$ | 8 | 1169.703 |
| v10e20_2 | 10 | 20 | 500 | $1.05 \mathrm{E}+6$ | 81.4761 | 80.7692 | $0.868 \%$ | $1 \%$ | 9 | 2725.251 |
| v10e20_3 | 10 | 20 | 500 | $1.05 \mathrm{E}+6$ | 81.6981 | 81.3323 | $0.448 \%$ | $1 \%$ | 10 | 1713.125 |
| v10e20_4 | 10 | 20 | 500 | $1.05 \mathrm{E}+6$ | 78.4703 | 78.1662 | $0.388 \%$ | $1 \%$ | 10 | 3164.391 |
| v10e20_5 | 10 | 20 | 500 | $1.05 \mathrm{E}+6$ | 77.3371 | 77.3371 | $0 \%$ | $1 \%$ | 10 | 4028.000 |
| v12e25_A | 12 | 25 | 300 | $3.36 \mathrm{E}+7$ | 75.5378 | 74.9683 | $0.754 \%$ | $1 \%$ | 8 | 2528.266 |
| v12e25_B | 12 | 25 | 300 | $3.36 \mathrm{E}+7$ | 52.4450 | 52.1022 | $0.654 \%$ | $1 \%$ | 12 | 3133.67 |
| v12e25_C | 12 | 25 | 300 | $3.36 \mathrm{E}+7$ | 70.3214 | 70.1937 | $0.182 \%$ | $1 \%$ | 11 | 1882.297 |
| v12e25_D | 12 | 25 | 300 | $3.36 \mathrm{E}+7$ | 43.1088 | 42.931 | $0.412 \%$ | $1 \%$ | 8 | 810.234 |
| v13e30_1 | 13 | 30 | 200 | $1.07 \mathrm{E}+9$ | 32.4429 | 32.3264 | $0.359 \%$ | $1 \%$ | 9 | 516.454 |
| v13e30_2 | 13 | 30 | 200 | $1.07 \mathrm{E}+9$ | 38.8679 | 38.6852 | $0.470 \%$ | $1 \%$ | 11 | 6332.030 |
| v13e30_3 | 13 | 30 | 200 | $1.07 \mathrm{E}+9$ | 32.4177 | 32.4131 | $0.014 \%$ | $1 \%$ | 7 | 1086.485 |
| v13e30_4 | 13 | 30 | 200 | $1.07 \mathrm{E}+9$ | 33.1784 | 32.9004 | $0.838 \%$ | $1 \%$ | 9 | 1095.703 |
| v13e30_5 | 13 | 30 | 200 | $1.07 \mathrm{E}+9$ | 34.4604 | 34.1707 | $0.841 \%$ | $1 \%$ | 9 | 3457.325 |
| v16e40_1 | 16 | 40 | 200 | $1.10 \mathrm{E}+12$ | 19.5682 | 19.4753 | $0.475 \%$ | $1 \%$ | 7 | 4367.515 |

